

# Nilpotency of square matrices with non-negative elements

Riccardo Bresciani\*, Mario Poletti\*\*, Andrew Butterfield\*

\*: Foundations and Methods Group, Trinity College Dublin & Lero@TCD — {bresciar, butrfield}@scss.tcd.ie

\*\* : Dipartimento di Matematica Applicata, Università di Pisa — mario.poletti@gmail.com



**Abstract** We give two equivalent decision procedures for nilpotency of square matrices with non-negative elements, which do not involve computation of the matrix eigenvalues. To the best of our knowledge these procedures have not been formalised elsewhere: we have decided to produce this brief technical report as a future reference.

A matrix is said to be *nilpotent of order  $k$*  if  $k$  is the least natural number such that:

$$M^k = 0$$

We aim at a decision procedure to enable us to decide by visual inspection if a matrix  $M \in \mathbb{R}^{n \times n}$  is nilpotent. In particular we are interested in matrices characterised by non-negative elements, and in this case we are sure that nilpotency can be inferred from considerations on the positions of null elements.

The most trivial observations are the following:

- if there is not at least one null column in  $M$ , then it is not nilpotent;
- if all columns of  $M$  are null (*i.e.*  $M$  is the null matrix), then it is nilpotent, and its nilpotency order is 1.

In the above cases  $M$  is said to be *terminal*, as its nilpotency is trivially decidable.

We are now going to describe informally how to understand where the other null elements should be so that a non-terminal matrix  $M$  is nilpotent.

If  $\text{Nilp}(M)$  denotes the nilpotency order of  $M$ , we have that:

$$\forall \underline{v} \in \mathbb{R}^n \bullet M^{\text{Nilp}(M)} \underline{v} = \underline{0}$$

If we imagine that  $\underline{v} = (\mu_1, \mu_2, \dots, \mu_n)^T$  represents a distribution of mass in  $n$  different places, then  $M$  can describe some event that leads to a new mass distribution, that depends on the previous one. The new distribution  $\underline{v}' = (\mu'_1, \mu'_2, \dots, \mu'_n)^T$  is such that:

$$\mu'_i = m_{i1}\mu_1 + m_{i2}\mu_2 + \dots + m_{in}\mu_n$$

and hence we can see this in the following way:

- if  $m_{ij} = 1$ , a copy of the mass in the  $j$ -th place is moved to the new  $i$ -th place:  $\mu'_i = \mu_j$ ;
- if  $m_{ij} = 0$ , a copy of the mass in the  $j$ -th place is trashed into the bin:  $\mu'_i = 0$ ;
- if  $m_{ij} < 1$ , a copy of the mass in the  $j$ -th place is shrunk and moved to the new  $i$ -th place and part is trashed into the bin:  $\mu'_i = m_{ij}\mu_j < \mu_j$ ;

\*The present work has emanated from research supported by Science Foundation Ireland grant 08/RFP/CMS1277 and, in part, by Science Foundation Ireland grant 03/CE2/I303.1 to Lero — the Irish Software Engineering Research Centre.

Nilpotent matrix

Null column

Null matrix

Terminal matrix

Informal description

- if  $m_{ij} > 1$ , a copy of the mass in the  $j$ -th place is enlarged and moved to the new  $i$ -th place:  
 $\mu'_i = m_{ij}\mu_j > \mu_j$ .

The subsequent application of the matrix  $M$  causes the mass to “follow” a certain path from its initial position to some other position: if every path leads to the bin, then we are happy to say that  $M$  is nilpotent.

We can describe the path in quite fine detail, as:

- if some  $k$ -th column is null, then all copies of the mass originally in the  $k$ -th place are thrown away after the first application of  $M$ ;
- if some  $h$ -th column is null with the exception of the element  $m_{kh}$  ( $k$  as above), then all copies of the mass originally in the  $h$ -th place are thrown away in two steps;
- if some  $g$ -th column is null with the exception of some elements among  $m_{hg}$  and  $m_{kg}$  ( $h, k$  as above), then all copies of the mass originally in the  $g$ -th are thrown away in at most three steps;
- ...

In the next section we formalise this procedure, we then provide a different procedure (with a proof), and finally provide an equivalence proof of the two procedures.

## 1 First procedure

Let  $M$  be non-terminal; given a succession  $Z_1, Z_2, \dots$  of sets let  $T_i$  be the set given by the union of the first  $i$  sets in the succession:

$$T_i \triangleq \bigcup_{1 \leq j \leq i} Z_j$$

Let  $Z_1, Z_2, \dots, Z_{(k+1)}$  be the succession such that:

- $Z_1$  is the set of the indices of all null columns of  $M$ ;
- if  $Z_i \neq \emptyset$ , then  $Z_{i+1}$  is the set of indices identifying the non-null columns of  $M$  such that the only non-null elements are those with row index contained in  $T_i$  or in a subset thereof;
- $Z_{(k+1)} = \emptyset$ .

Then  $M$  is nilpotent if and only if  $T_k = \{1, 2, \dots, n\}$ ; moreover in this case it is  $\text{Nilp}(M) = k$ .

If we look at the succession  $Z_1, Z_2, \dots, Z_{(k+1)}$  from the perspective used in the informal description, we have that:

- $Z_1$  indicates all places from where the mass is thrown away;
- more in general,  $Z_i$  indicates what are the places from where the mass is thrown away in  $i$  steps;
- therefore  $T_i$  indicates what are the places from where the mass is thrown away in at most  $i$  steps.

## 2 Second procedure

Let  $M$  be non-terminal;  $M^R$  is the *reduction* (or *reduced matrix*) of  $M$ , if it can be obtained from  $M$  by removing all rows and columns with indices contained in  $Z_1$ , where  $Z_1$  is the set of all indices of null columns (as in procedure 1).

Let  $M_1, M_2, \dots, M_k$  be the succession such that:

- $M_1 = M$ ;
- if  $M_i$  is not terminal, then  $M_{i+1} = M_i^R$ ;

Set  $T_i$ 

Procedure 1

Reduced matrix

Succession

- $M_k$  is terminal.

Then  $M$  is nilpotent if and only if  $M_k$  is nilpotent; moreover in this case all the matrices in the succession are nilpotent, and their nilpotency orders verify the relations

$$\text{Nilp}(M_k) = 1 \quad \text{Nilp}(M_{(i-1)}) = \text{Nilp}(M_i) + 1$$

and therefore

$$\text{Nilp}(M) = k$$

The procedure to decide on the nilpotency of  $M$  is therefore to start with the matrix  $M$  and subsequently cross out all columns containing exclusively null elements along with the rows with the same index: by iterating the process we eventually end up with a reduction  $M_k$  of  $M$  which contains only non-null columns or only null columns, and therefore the nilpotency of  $M_k$  is trivially decidable.

Procedure 2

## 3 Proofs

### 3.1 Proof of procedure 2

Let  $P$  be a permutation matrix such that:

$$PMP^{-1} = \begin{pmatrix} 0_{r \times r} & H_{r \times s} \\ 0_{s \times r} & M^R \end{pmatrix}$$

which contains  $r$  null columns, with indices less or equal than  $r$ , and  $s = n - r$  non-null columns, with indices greater than  $r$ .

We then have that:

$$PM^hP^{-1} = (PMP^{-1})^h = \begin{pmatrix} 0_{r \times r} & H_{r \times s}(M^R)^{(h-1)} \\ 0_{s \times r} & M^R(M^R)^{(h-1)} \end{pmatrix}$$

We also have that  $M^h = 0 \Leftrightarrow PM^hP^{-1} = 0$ , and therefore:

$$M^h = 0 \Leftrightarrow \begin{pmatrix} H_{r \times s} \\ M^R \end{pmatrix} (M^R)^{(h-1)} = 0$$

which can be true if and only if  $(M^R)^{(h-1)} = 0$ , given that in the matrix

$$\begin{pmatrix} H_{r \times s} \\ M^R \end{pmatrix}$$

all columns are non-null, and all elements are non-negative.

We can therefore conclude that:

- $M$  is nilpotent if and only if  $M^R$  is nilpotent;
- $\text{Nilp}(M) = \text{Nilp}(M^R) + 1$ .

It should be noted that removing the assumption  $m_{ij} \geq 0$  allows us to conclude that in general

$$\text{Nilp}(M^R) \leq \text{Nilp}(M) \leq \text{Nilp}(M^R) + 1$$

**Observation** Given that  $M^h \rightarrow 0 \Leftrightarrow PM^hP^{-1} \rightarrow 0$ , we can show with a similar proof that

$$M^h \rightarrow 0 \Leftrightarrow (M^R)^h \rightarrow 0$$

---

### 3.2 Proof of equivalence of the procedures

The set  $Z_i$  of procedure 1, contains the indices of the columns and rows which are crossed out at the  $(i + 1)$ -th step of procedure 2; hence the equivalence of the procedures is proved by the following remarks:

- for  $i = 2, \dots, k$  the matrix  $M_i$  of procedure 2 can be obtained from  $M$  by removing all rows and columns with indices contained in  $T_{i-1}$  of procedure 1;
- the matrix  $M_k$  is terminal and null if and only if  $Z_k \neq \emptyset$  and  $T_k = \{1, \dots, n\}$ ; moreover in this case it is  $Z_{k+1} = \emptyset$ ;
- the matrix  $M_k$  is terminal and non null if and only if  $Z_k = \emptyset$  and  $T_{(k-1)} \neq \{1, \dots, n\}$ .